Effective DB2 Object Monitoring Using The DB2 Catalog

By Craig S. Mullins

n order to maintain efficient production DB2-based systems, it is necessary to monitor periodically the DB2 objects comprising those systems. Monitoring is an essential component of postimplementation duties because the production environment is dynamic. Fluctuations in business activity, errors in logical and/or physical design or lack of communication can all cause a system to perform inadequately. An effective strategy for monitoring DB2 objects in the production environment will catch and forestall potential problems before they adversely impact performance. The job of monitoring DB2 objects usually is performed by a database administrator, performance analyst or system administrator.

An effective monitoring strategy will include scheduled performance monitor jobs, regular monitor runs from all online environments where DB2 transactions execute (that is, CICS, IMS/DC, TSO) and regularly-scheduled reports from the DB2 catalog. This article will detail a strategy for accomplishing the latter objective.

By querying the DB2 catalog on a consistent basis, many possible hazards can be analyzed and acted on in a proactive mode. This article describes QMF queries which access the DB2 catalog. An analysis accompanies each query highlighting some of the potential problems that can be trapped by reviewing the output from each query.

To implement this strategy, certain basic assumptions have been made. It is assumed that all application plans are bound with the EXPLAIN(YES) option and that each application has its own PLAN_TA-

BLE for the storage of the EXPLAIN results. It is also assumed that scheduled production STOSPACE and RUNSTATS jobs are executed and that plans are rebound whenever RUNSTATS has been executed. This is necessary to maintain current statistical information about all DB2 objects for an application. It is also assumed that RUNSTATS has been performed on the DB2 catalog tablespaces. This will allow the optimal performance of these queries.

It will also be useful to have a report of each PLAN_TABLE for each application. This will allow for cross-checking the DB2 catalog against application plan and optimizer path selection information. Run the query in Figure 1 for each PLAN_TABLE in order to produce these reports. Note: It is crucial for the TIME-STAMP column to be in descending order. This will cause the EXPLAIN data in the report to be sorted in order from the most recent to the oldest for each query in the PLAN TABLE.

This is important if the PLAN_TA-BLE(s) being used are not purged. The DB2 Application Programming Guide contains information about EXPLAIN and the PLAN_TABLE.

The queries and forms presented in this article were developed using QMF and are run weekly using a batch QMF job. This is easier than submitting the queries from QMF weekly or through SPUFI. Simply build batch QMF JCL, incorporate all these queries and forms into a proc and run the proc.

If these basic assumptions and conditions are not met, it is still possible to implement the queries. To accomplish

this, changes may need to be made to the SQL or to the process by which they are submitted. In addition, if STOSPACE, RUNSTATS and EXPLAIN data is not current, the queries may not provide accurate information. With these basic caveats in mind, a description of each catalog query follows.

The Object Listing Queries

• Table Listing — Lists table columns • Index Listing — Lists index columns

In order to perform database and system administration functions for DB2, often it is necessary to identify quickly object dependencies. For example, a DBA is confronted with the need to analyze a poorly-performing query. He has the query and a report of the EXPLAIN for the query. Missing is a listing of available indexes and candidate columns for creating indexes. The object listing queries (see Figures 2 and 3) provide this and more.

By viewing the output from these two queries a hierarchy of DB2 objects can be easily ascertained (indexes within tables within tablespaces within databases). The output from these queries is superb for navigational purposes. It is easy to get lost in a flood of production objects. By periodically running these queries and saving the output, a DBA can have an upto-date profile of the environment established in each DB2 subsystem that he must monitor.

The WHERE T.DBNAME IN clause in each of these queries is optional. This clause is used to produce reports only for those databases which need to be monitored. It is usually desirable to eliminate the sample database, the DB2 catalog da-

PLAN TABLE Query Query QUERYNO, QBLOCKNO, APPLNAME, PROGNAME, PLANNO, METHOD, CREATOR, TNAME, TABNO, ACCESSTYPE, SELECT MATCHCOLS, ACCESSNAME, INDEXONLY, SORTN_UNIQ, SORTN_JOIN, SORTN_ORDERBY, SORTN_GROUPBY, SORTC_UNIQ, SORTC_JOIN, SORTC_ORDERBY, SORTC_GROUPBY, TSLOCKMODE, TIMESTAMP, PREFETCH, COLUMN_FN_EVAL, MIXOPSEQ FROM ownerid.PLAN TABLE ORDER BY APPI NAME **PROGNAME** TIMESTAMP DESC. QUERYNO. PLANNO Form FORM, COLUMNS userid.PLANTBF Total Width Of Report Columns: 148 NUM COLUMN HEADING USAGE INDENT WIDTH EDIT QBN APPLNAME PROGNAME BREAK1 PN METHOD CREATOR C TNAME c TN 9 TN 10 AT 11 MCOL 12 XNAME 13 IX_ONLY 14 S_U 15 S_J 16 S_6 C 10 13 00000000 15 16 17 S O 18 C U 19 C J 20 C 6 21 C O 22 TSL 18 19 21 22 23 24 TIMESTAMP CC 24 COL EVAL 25 MULT_INDX

tabase and any other extraneous databases (that is, DBEDIT, QMF and so on). This elimination is optional as a DBA may wish to monitor everything known to DB2.

Although the primary purpose of these two queries is navigational, they can also aid in problem determination and performance tuning. For example, note the following query:

SELECT A.COL1, A.COL2, B.COL3 FROM TABLE1 A, TABLE2 B WHERE A.COL1 = B.COL4;

If this query is not performing properly, the DBA needs to know the column types and lengths for COL1 in TABLE1 and COL4 in TABLE2. The type and length for both columns should be the same. If they are not then DB2 is performing a data conversion to make the comparison. This will negatively impact performance. Should the data type and length be consistent then find what indexes are defined (if any) on these columns and analyze the EXPLAIN output. Other data such as the key ranges and cardinality may also be significant, whether an index is clustered or not (these items will influence the optimizer's choice of access path) and the number of tables in a tablespace (this may cause performance degradation for non-segmented tablespaces). All of this information can be obtained from these reports.

This is only one level of DB2 performance tuning. The next level will delve

FIGURE 2

		Table Listing	Query				
Que	ry						
	SELECT	T.DBNAME, T.TSNAME, T.CREATOR, C.NAME, C.COLTYPE, C.LENGTH, C.S C.DEFAULT, C.COLCARD, C.HIGH2KE C.FLDPROC	CALE, C.	VULLS,	PDATES,		
	FROM	SYSIBM.SYSCOLUMNS C.					
		SYSIBM.SYSTABLES T					
	WHERE		T.NAME	E = C.TBN	IAME)		
	AND	T.TYPE = 'T'					
	AND	T.DBNAME IN ('dbname1', 'dbname2',	dbnamex)			
	ORDER B						
		T.DBNAME,					
		T.TSNAME,					
		T.NAME, C.COLNO					
		C.COLNO					
Form	1						
Congression of							
FORM	COLUMNS		userid.LIS	SHBLE			
Total V	Width Of Re	port Columns: 135					
NIIM	COLUMN	HEADING	USAGE	INDENT	WIDTH	FDIT	SEQ
	COLOMIT	педина				Control of the last	
1			BREAK1		8	C	1
	TABLE S		BREAK2		8	C	2
	TABLE_C	REATOR	BREAK3	1	8	C	3
	TABLE		BREAK3	1	18	C	4
	COL_NO			1	3	L	5
	COLUMN				18	C	7
	COLUMN			1	6	L	8
	COLUMN_ SCALE	LENGIH			6	i	9
	NU LL				2	C	10
	DF LT				2	C	11
	COL CAP	ND.			8	Ľ	12
	HIGH2 KI			1	8	C	13
	LOW2 KE			1	8	C	14
	UPDT			1	4	C	15
	FLD PRO	C			4	č	16
10	. Lo no			STATE OF THE PARTY.			

FIGURE 3

			ndex Listing Q	uery				
Que	ry							
	SELECT	I.UNIQUERULE, I.C.	ME, T.CREATOR, T.N. LUSTERING, I.CLUST FULLKEYCARD, I.N. Q. K.COLNAME, K.OF K.	ERED, LEAF, I.N	LEVELS,			
	A CONTRACTOR OF THE PARTY OF TH	SYSIBM.SYSTABLE	S T,					
		SYSIBM.SYSINDEX						
	WHERE		CREATOR AND I.TBI					
	AND		ame1', 'dbname2', 'db					
	ORDER E		unici, donamez, de	Jiloinea ,	Control of the last of the las			
		T.DBNAME,						
		T.TSNAME,						
		T.NAME,						
		I.NAME, K.COLSEQ						
		K.COLSEU						
Forn	n							
	-							
	.COLUMN	S	u	serid.LIS	STINDF			
FORM		S eport Columns: 136	u	serid.LIS	STINDF			
FORM Total	Width Of R			serid.LIS		WIDTH	EDIT	SEQ
FORM Total	Width Of R	eport Columns: 136	<u>u</u> B	ISAGE IREAK1	INDENT 1	8	C	1
FORM Total 1 NUM 1 2	COLUMN DATABLE	eport Columns: 136 HEADING ASE BPACE	U B B	SAGE REAK1	INDENT 1 1	8 8	C	1 2
FORM Total 1 NUM 1 2 3	COLUMN DATABLE TABLE	eport Columns: 136 HEADING ASE BPACE	U B B B	SAGE REAK1 REAK2 REAK3	INDENT 1 1	8 8 8	CCC	1 2 3
FORM Total 1 NUM 1 2 3 4	COLUMN DATABLE TABLE TABLE TABLE	eport Columns: 136 HEADING ASE BPACE	U 8 8 8	REAK1 REAK2 REAK3	INDENT 1 1 1 1 1	8 8 8 18	0000	1 2 3 4
FORM Total 1 NUM 1 2 3 4 5	COLUMN _DATAB/ TABLE_S TABLE_INDEX	eport Columns: 136 HEADING ASE BPACE	U 8 8 8	SAGE REAK1 REAK2 REAK3	INDENT 1 1 1 1 1	8 8 8 18 18	00000	1 2 3 4 5
FORM Total 1 NUM 1 2 3 4 5 6	COLUMN _DATAB/ TABLE_C TABLE_INDEX U_Q	eport Columns: 136 HEADING ASE BPACE	U 8 8 8	REAK1 REAK2 REAK3	INDENT 1 1 1 1 1	8 8 8 18 18	000000	1 2 3 4 5 6
Total 1 NUM 1 2 3 4 5 6 7	COLUMN DATABLE TABLE TABLE INDEX U C L	eport Columns: 136 HEADING ASE BPACE	U 8 8 8	REAK1 REAK2 REAK3	INDENT 1 1 1 1 1	8 8 8 18 18 1 1	0000000	1 2 3 4 5 6 7
Total 1 NUM 1 2 3 4 5 6 7 8	COLUMN DATABLE TABLE TABLE INDEX U_Q C_L C_D	eport Columns: 136 HEADING ASE BPACE CREATOR	U 8 8 8	REAK1 REAK2 REAK3	INDENT 1 1 1 1 1	8 8 8 18 18 1 1	000000	2 3 4 5 6
Total 1 1 2 3 4 5 6 7 8 9	COLUMN DATABLE TABLE TABLE INDEX U Q C_L C_D 1ST KEY	eport Columns: 136 HEADING ASE SPACE CREATOR	U 8 8 8	REAK1 REAK2 REAK3	INDENT 1 1 1 1 1	8 8 8 18 18 1 1	00000000	1 2 3 4 5 6 7 8
Total 1 1 2 3 4 5 6 7 8 9 10	COLUMN DATABLE STABLE SINDEX U_Q C_L C_D 1ST KEY FULL KE	eport Columns: 136 HEADING ASE SPACE CREATOR CARDINAL Y CARDINAL	U 8 8 8	REAK1 REAK2 REAK3	INDENT 1 1 1 1 1	8 8 8 18 18 1 1 1 1 8	00000000	1 2 3 4 5 6 7 8 9
Total 1 1 2 3 4 5 6 7 8 9	COLUMN DATABLE TABLE TABLE INDEX U_Q C_L C_D 1ST KEY FULL KE NO OF 1	eport Columns: 136 HEADING ASE SPACE CREATOR CARDINAL Y CARDINAL EAFS	U 8 8 8	REAK1 REAK2 REAK3	INDENT 1 1 1 1 1	8 8 8 18 18 1 1 1 1 8 8 6 6	00000001111	1 2 3 4 5 6 7 8 9
Total 1 NUM 1 2 3 4 5 6 7 8 9 10 11	COLUMN DATABLE STABLE SINDEX U_Q C_L C_D 1ST KEY FULL KE	eport Columns: 136 HEADING ASE SPACE CREATOR CARDINAL Y CARDINAL EAFS	U 8 8 8	REAK1 REAK2 REAK3	INDENT 1 1 1 1 1	8 8 8 18 18 1 1 1 1 8 8 8	0000000011110	1 2 3 4 5 6 7 8 9 10
Total 1 NUM 1 2 3 4 5 6 7 8 9 10 11 12	COLUMN DATABLE TABLE INDEX U_Q C_L C_D 1ST KEY FULL KE NO OF LI NO OF LI	eport Columns: 136 HEADING ASE BPACE CREATOR CARDINAL Y CARDINAL EAFS EVELS	U 8 8 8	REAK1 REAK2 REAK3	INDENT 1 1 1 1 1	8 8 8 18 18 1 1 1 1 8 8 6 6 6	00000000111101	1 2 3 4 5 6 7 8 9 10 11
Total 1 NUM 1 2 3 4 5 6 7 8 9 10 11 12 13	COLUMN DATAB/ TABLE TABLE INDEX U_Q C_L C_D IST KEY FULL KE NO OF L NO OF L ISOBID	eport Columns: 136 HEADING ASE SPACE CREATOR CARDINAL Y CARDINAL EAFS EVELS	U 8 8 8	REAK1 REAK2 REAK3	INDENT 1 1 1 1 1	8 8 8 118 118 1 1 1 1 8 8 6 6	0000000011110	1 2 3 4 5 6 7 8 9 10 11 12 13

deeper into the physical characteristics of DB2 objects.

The Physical Statistics Queries

- Tablespace Physical Statistics
- Index Space Physical Statistics

Quite often it will be necessary to trace a performance problem within a DB2 query to the physical level. Characteristics at the physical level are determined when DB2 objects are defined. The focus will be on tablespaces and index spaces as these two objects require that a physical data set be created to support them. Many different options need to be chosen when a DB2 is created. If poor choices are made, performance will be unsatisfactorily impacted. The physical statistics queries can be used to monitor and tune these options.

The Tablespace Physical Statistics Query (see Figure 4) provides a listing of tablespaces within the database and lists all the pertinent physical detail associated with each tablespace. The Index Physical Statistics Query (see Figure 5) provides a report of all indexes grouped by owner with the physical criteria supporting each index. These reports are invaluable tools for diagnosing performance problems when they happen and frequently for catching problems before they occur.

Both reports in this section show the CLOSE RULE associated with the ta-

blespace or index space. Always monitor this rule for both. A CLOSE RULE of Y indicates that every time an object is accessed, a VSAM open and close will be performed by the system. The performance of any query which accesses an object defined this way will be impeded. A CLOSE RULE of N will only perform the VSAM open the first time the object is accessed. It will then remain open until DB2 is shut down. This will increase the performance of the query. It does, however, add the overhead associated with keeping a data set open. This overhead should be minimal and is usually preferred to having a slow-running query.

Each tablespace and index must be reviewed based on its desired usage to determine the CLOSE RULE for it. Certain objects which are accessed infrequently or only once per day will not need to remain open. As a basic rule, define tablespaces as CLOSE NO unless a good reason is provided otherwise. Whenever a query is found to be causing performance problems, always examine the CLOSE RULE for each tablespace and index involved in the query.

These reports are also useful in determining frequency of reorganization. By monitoring PCT Dropped, Far Off and Near Off Rows (in both the tablespace and index spaces), Leaf Distance and Cluster Ratio it can be determined whether

to increase or decrease the frequency of running a reorganization. See Table 1 for an analysis of the impact of this information on reorganization frequency.

It is also useful to analyze the table-space and index space usage. The ability to monitor SPACE USED% for a table-space is of particular importance. Efficient DASD space usage is important to maintain an optimal operating environment for DB2 objects. When the SPACE USED% consistently remains below 75 percent for an extended period of time, the PRIQTY space for the tablespace should be decreased and the tablespace should be reorganized. If this percentage

TABLE 1

Reorganization Indicators					
Column	Object	Impact			
PCT Dropped	TS	++++			
Near Off Rows	TS	+			
Far Off Rows	TS	++++			
Cluster Ratio	Index				
Near Off Rows	Index	+			
Far Off Rows	Index	++++			
Leaf Distance	Index	+++			

How To Read The Above Chart

- A'+' indicates REORG more frequently when this number is large
- A'-' indicates REORG more frequently when this number is small

The greater the number of occurrences the more urgent the need to REORG

FIGURE 4

FIGURE

Index Physical Statistics Query

		Tablespace Physic	cai Statistic	s Query			
Que	ry						
	SELECT	T.DBNAME, T.NAME, T.IMPLICIT, T.ERASERULE, T.CLOSERULE, T.NTABLES, T.NACTIVE, P.CARC P.NEARINDREF, P.PERCACTIVE, P.FREEPAGE, P.PCTFREE, P.STO T.NACTIVE*100*T.PGSIZE/T.SPA	T.PARTITIONS, T D, P.FARINDREF, P.PERCDROP, DRNAME, P.VCAT				
	FROM	SYSIBM.SYSTABLESPACE T, SYSIBM.SYSTABLEPART P					
	WHERE AND ORDER E	T.NAME = P.TSNAME T.DBNAME = P.DBNAME					
	OHDER	T.DBNAME, T.NAME					
Forn	n						
FORM	COLUMN		userid.PH	YTABLE			
Total \	Width Of R	eport Columns: 150					
NUM	COLUMN	HEADING	USAGE	INDENT	WIDTH	EDIT	SEQ
NUM 1	COLUMN			INDENT 1	WIDTH 8	EDIT C	SEQ 1
1 2	A STATE OF THE PARTY OF THE PAR	ISE		1	8 8	C	1 2
1 2 3	_DATABATABLE_S	ISE PACE	BREAK1	1	8 8 4	CCCC	1 2 3
1 2 3 4	DATABA TABLE_S IMPL LOCK_R	ISE PACE	BREAK1	1 1	8 8 4 4	CCCC	1 2 3 4
1 2 3 4 5	DATABA TABLE_S IMPL LOCK_RI	ISE PACE	BREAK1	1 1 1 1 1	8 8 4 4 1	00000	1 2 3 4 5
1 2 3 4 5 6	DATABA TABLE_S IMPL LOCK_RI E_R C_R	ISE PACE	BREAK1	1 1 1 1 1 1 1	8 8 4 4 1	000000	1 2 3 4 5 6
1 2 3 4 5 6 7	DATABA TABLE_S IMPL LOCK_RI E_R C_R PARTS	ISE IPACE ULE	BREAK1	1 1 1 1 1 1 1 1	8 8 4 4 1 1 5	CCCCCCL	1 2 3 4 5 6 7
1 2 3 4 5 6 7 8	DATABA TABLE_S IMPL LOCK_RI E_R C_R PARTS SEG SIZ	ISE PACE JULE	BREAK1	1 1 1 1 1 1 1 1	8 8 4 4 1 1 5 4	CCCCCCLL	1 2 3 4 5 6 7 8
1 2 3 4 5 6 7 8 9	DATABA TABLE_S IMPL LOCK_RI E_R C_R PARTS SEG_SIZ NO OF_T	SE PACE ULE E ABLES	BREAK1	1 1 1 1 1 1 1 1	8 8 4 4 1 1 5 4 6	000000111	1 2 3 4 5 6 7 8
1 2 3 4 5 6 7 8 9	DATABA TABLE S IMPL LOCK RI E_R C_R PARTS SEG_SIZ NO OF_T NO OF_F	ISE IPACE JLE E ABLES AGES	BREAK1	1 1 1 1 1 1 1 1	8 8 4 4 1 1 5 4 6 8	0000001111	1 2 3 4 5 6 7 8 9
1 2 3 4 5 6 7 8 9 10	DATABA TABLE_S IMPL LOCK_RI E_R C_R PARTS SEG_SIZ NO OF_T NO OF_F	SE PACE ULE E ABLES AGES OWS	BREAK1	1 1 1 1 1 1 1 1	8 8 4 4 1 1 5 4 6 8	00000011111	1 2 3 4 5 6 7 8 9 10
1 2 3 4 5 6 7 8 9 10 11	DATABA TABLE_S IMPL LOCK_R E_R C_R PARTS SEG_SIZ NO OF_T NO OF_F FAR OFF	ISE PACE ULE E ABLES AGES OWS ROWS	BREAK1	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	8 8 4 4 1 1 5 4 6 8 11	000000111111	1 2 3 4 5 6 7 8 9 10 11 12
1 2 3 4 5 6 7 8 9 10 11 12	DATABA TABLE_S IMPL LOCK_RI E_R C_R PARTS SEG_SIZ NO OF_T NO OF_F NO OF_F NO OF_F NEAR OF	ISE IPACE ULE E ABLES AGES IOWS ROWS F ROWS	BREAK1	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	8 8 4 4 1 1 5 4 6 8 11 11	0000001111111	1 2 3 4 5 6 7 8 9 10 11 12 13
1 2 3 4 5 6 7 8 9 10 11 12 13 14	DATABA TABLE_S IMPL LOCK_RI E_R C_R PARTS SEG_SIZ NO OF_T NO OF_F NO OF_F FAR OFF NEAR OF	ISE PACE ULE E ABLES AGES OWS ROWS F ROWS COTIVE	BREAK1	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	8 8 4 4 1 1 5 4 6 8 11 11 11 7	00000011111111	1 2 3 4 5 6 7 8 9 10 11 12 13 14
1 2 3 4 5 6 7 8 9 10 11 11 12 13 14 15	DATABA TABLE S IMPL LOCK RI E_R C_R PARTS SEG_SIZ NO OF_T NO OF_R NO OF_R FAR OFF NEAR OF	ISE IPACE JLE ABLES AGES IOWS F. ROWS CTIVE INFOPPED	BREAK1	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	8 8 4 4 1 1 5 4 6 8 11 11 17 7	0000001111111111	1 2 3 4 5 6 7 8 9 10 11 12 13 14 15
1 2 3 4 5 6 7 8 9 10 11 11 12 13 14 15 16	DATABA TABLE S IMPL LOCK RI E R C R PARTS SEG SIZ NO OF T NO OF P NO OF P NO OF P NEAR OF PCT A PCT A PCT A	SE PACE ULE E ABLES AGES OWS ROWS F, ROWS CTIVE GROPPED GE	BREAK1	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	8 8 4 4 1 1 5 4 6 8 111 111 7 7	0000001111111111	1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16
1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17	DATABLE SIMPL LOCK RIE R PARTS SEG SIZ NO OF T NO OF FAR OFF NO ACT OF PCT ACT OF PCT ACT OF PCT ACT OF FREE PARTS	ISE IPACE JLE ABLES AGES OWS ROWS F. ROWS CÜTVE IROPPED GE	BREAK1	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	8 8 4 4 1 1 5 4 6 8 111 117 7 7 4 4	000000111111111111	1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17
2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18	DATABA TABLE SIMPL LOCK_RIES R E.R. C_R PARTS SEG SIZ NO OF TOO OF FAR OF FAR OF FAR OF FAR OF PCT L PCT L FREE PAPCT FRESTORO	SE PACE ULE E ABLES AGES OWS ROWS F, ROWS CTIVE PROPPED RGE E E	BREAK1	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	8 8 4 4 1 1 5 4 6 8 11 11 7 7 7 7	0000000111111111110	1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18
1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17	DATABLE SIMPL LOCK RIE R PARTS SEG SIZ NO OF T NO OF FAR OFF NO ACT OF PCT ACT OF PCT ACT OF PCT ACT OF FREE PARTS	EE PACE ULE E ABLES AGES OWS ROWS CTIVE PROPPED GGE EE UP	BREAK1	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	8 8 4 4 1 1 5 4 6 8 111 117 7 7 4 4	000000111111111111	1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17

Que	· y						
	FROM WHERE AND ORDER E	LCREATOR, LCREATEDBY, LNA LCLUSTERING, LCLUSTERED, I LFIRSTKEYCARD, LFULLKEYCA 4096/LPGSIZE, LERASERULE, I. PFAROFFPOS, PLEAFDIST, PN PPCTFREE SYSIBM.SYSINDEXES I, SYSIBM.SYSINDEXPART P LNAME = PLXNAME LCREATOR = PLXCREATOR LY LCREATOR, LNAME	CLUSTERRATIO ARD, I.NLEAF, I.N CLOSERULE, P.	LEVELS, CARD,	Ξ,		
Forn	0						
	.COLUMN:		userid.PH	IYINDXF			
Total \	Width Of R	eport Columns: 143					
NUM	COLUMN	HEADING	USAGE	INDENT	WIDTH	EDIT	SEQ
1	INDEX_C	WNER	BREAK1	1	8	C	1
2	INDEX			1	18	C	2
3	U_Q			1	1	C	3
4	C_L			1	1	C	4
5	C_D			1	1	C	5
6	CLSTR_F			1	5	L	6
7		Y_CARDINALITY		1	11	L	7
8		Y_CARDINALITY		1	11	L	8
9	NO OF_LEAFS			1	6	L	9
10	NO OF_L			1	6	L	10
11	SUB_PAG	BES		1	5	L	11
12	E_R			1	1	C	12
13	C_R			1	1	C	13
14		DWS_REF (CARD)		1	11	L	14
15		FAROFF POS		1	11	L	15
16	LEAF_DI			1	8	L	16
17		NEAROFF POS		1	11	L	17
18	PCT FRE			1	4	L	18
19							

6 U **Tablespace Scan Query** Query SELECT E.APPLNAME, E.PROGNAME, E.QUERYNO, E.TNAME, T.NPAGES, E.TIMESTAMP, S.SEQNO, S.TEXT ownerid.PLAN_TABLE SYSIBM.SYSTABLES FROM SYSIBM.SYSSTMT S ACCESSTYPE = 'R' (T.NPAGES > 50 OR T.NPAGES < 0) AND (I.NPAGES > 50 OH I.NPAG T.NAME = E.TNAME T.CREATOR = E.CREATOR S.NAME = E.PROGNAME S.PLNAME = E.APPLNAME S.STMTNO = E.QUERYNO AND AND AND AND ORDER BY E.APPLNAME, **E.PROGNAME** E TIMESTAMP DESC S.SEQNO Form userid.SCANTSF FORM.COLUMNS Total Width Of Report Columns: 329 INDENT WIDTH EDIT SEQ NUM COLUMN HEADING C PLAN_NAME DBRM_NAME BREAK2 BREAK3 BREAK3 QRY NO TABLE_NAME NO OF_PAGES _TIME STAMP **BREAK3** 16

		Index Space	ce Scan Quer	у			
Que	<u>y</u>						
	SELECT	E.APPLNAME, E.PROGNAME, I.COLCOUNT, E.MATCHCOLS, S.SEQNO, S.TEXT,					
	FROM	ownerid.PLAN_TABLE E, SYSIBM.SYSINDEXES I, SYSIBM.SYSSTMT S					
	WHERE	E.ACCESSTYPE = 'I'					
	AND	N.LEAF > 100					
	AND	E.MATCHCOLS < I.COLCOUN	П				
	AND	I.NAME = E.ACCESSNAME					
	AND	I.CREATOR = E.ACCESSCRE	ATOR				
	AND	S.NAME = E.PROGNAME S.PLNAME = E.APPLNAME					
	AND	S.STMTNO = E.QUERYNO					
	ORDER E						
	OHDERL	E.APPLNAME,					
		E.PROGNAME,					
		E.TIMESTAMP DESC.					
		E.QUERYNO,					
		S.SEQNO					
Forn	1						
FORM	COLUMN	3	userid.SC	ANIXF			
Total \	Width Of F	eport Columns: 351					
		HEADING	USAGE	INDENT	WIDTH	EDIT	SEQ
NUM	COLUMN	TIEADING		-			
NUM 1	PLAN N		BREAK1		8	C	1
		AME	BREAK2	1	8	C	2
1 2 3	PLAN_N DBRM_N QRY_NC	AME IAME	BREAK2 BREAK3	1	8 7	C	2 3
1 2 3 4	PLAN_N DBRM_N QRY_NO INDEX	AME AME	BREAK2 BREAK3 BREAK3	1	8 7 18	CLC	2 3 4
1 2 3 4 5	PLAN_N DBRM_N QRY_NO INDEX NO OF_I	AME IAME LEAF PAGES	BREAK3 BREAK3 BREAK3	1 1 1 1	8 7 18 8	CLCL	2 3 4 5
1 2 3 4 5 6	PLAN_N DBRM_N QRY_NO INDEX NO OF_I COLUMN	AME IAME EAF PAGES S IN KEY	BREAK2 BREAK3 BREAK3 BREAK3 BREAK3	1 1 1 1	8 7 18 8 6	CLCLL	2 3 4 5 6
1 2 3 4 5 6 7	PLAN_N DBRM_N QRY_NO INDEX NO OF L COLUMN MATCHIN	AME IAME EAF PAGES S IN KEY IG COLUMNS	BREAK2 BREAK3 BREAK3 BREAK3 BREAK3 BREAK3	1 1 1 1 1 1	8 7 18 8	CLCLLL	2 3 4 5
1 2 3 4 5 6 7 8	PLAN_N DBRM_N QRY_NO INDEX NO OF L COLUMN MATCHIN INDX_OR	AME LAME LEAF PAGES S. IN KEY IG_COLUMNS	BREAK2 BREAK3 BREAK3 BREAK3 BREAK3	1 1 1 1	8 7 18 8 6 8	CLCLL	2 3 4 5 6 7
1 2 3 4 5 6 7	PLAN_N DBRM_N QRY_NO INDEX NO OF L COLUMN MATCHIN	AME IAME LEAF PAGES S. IN KEY IG. COLUMNS ILY	BREAK2 BREAK3 BREAK3 BREAK3 BREAK3 BREAK3 BREAK3	1 1 1 1 1 1 1 1	8 7 18 8 6 8 4	0 1 0 1 1 1 0	2 3 4 5 6 7 8

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is 100 percent and growth is expected, the PRIQTY space should be increased.

Other space considerations can be analyzed via monitoring PCT ACTIVE, FREE PAGE and PCT FREE. In conjunction with a data set allocation report, space can be reviewed and modified as necessary. As a general rule when PCT ACTIVE is low it may be necessary to redefine with a smaller PRIQTY (and/or SECQTY). Free space can also be changed as well. In any event it will be necessary to monitor these reports with the actual data set statistics. Also, remember that in order for changes-to-space characteristics to take affect, the tablespace being altered must be reorganized.

After this level of performance analysis has been exhausted it becomes necessary to broaden the scope of the tuning effort. This will involve analysis of application programs (plans) and may necessitate building new indexes or changing SQL in application queries.

The Scan Queries

- Tablespace Scans Greater Than 50 Pages
- Index Space Scans Greater Than 100 Pages

The scan queries (see Figures 6 and 7) produce reports that can detect many potential performance problems. By combining the DB2 catalog information with the output from EXPLAIN, a series of potential "problem queries" can be identified. These "problem queries" will be

grouped into two categories; tablespace scans and index space scans. In each instance, DB2 will scan data sets to satisfy the query. A tablespace scan will not use an index and will read every page in the tablespace. An index space scan will not necessarily read every index subpage but has that potential.

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These queries will probably be long-running. Do not execute them in parallel with heavy production DB2 processing or during the on-line DB2 transaction window. For the scan queries to operate efficiently, ensure that the PLAN_TABLE being used in each query does not contain excessive extraneous data. Strive to maintain only the most recent EXPLAIN data from production BIND jobs in the table. Also, only keep EXPLAIN information for plans which need to be monitored. Executing RUNSTATS on the PLAN_TABLES can also positively affect the performance of these queries.

The tablespace scan report will list queries scanning more than 50 pages and queries accessing tables without current RUNSTATS information. If the NO OF PAGES is -1 for any table then RUNSTATS has not been run. A RUNSTATS job should be executed as soon as possible followed by a rebind of any plan which uses this table. Everything else on this report should be monitored closely. For tables over the 50-page threshold, the performance impact is uncertain. The greater

the number of pages scanned, the greater the potential for performance problems.

The 50-page cutoff for inclusion on this report is somewhat arbitrary and may need to be redefined as the information returned is gauged. To monitor large tables only, the number may be increased. If a bufferpool is small (under 1000 buffers), the number may be reduced. Some potential numbers and their significance follow:

- 100 Increase this number to 100 (or larger) to monitor only those queries accessing large tables. This number will vary according to the definition of "large table."
- 20 For tables with 20 or more pages it is recommended that indexes be created in order to satisfy the predicates in the query. It is not always possible to create an index for every predicate however, so this is only a guideline. Various DB2 references recommend indexes be considered when the number of pages in a tablespace reaches five, six and/or 15. In practice though, 20 pages seems to be a good number.

The index space scan report will list all queries scanning more than 100 leaf pages where a match on the columns in the query is not a complete match on all index columns. As the number of MATCHING COLUMNS increases the possibility of performance problems decrease. The

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simple COBOL programs for the product to pay for itself," he adds.

For NIPSCO the initial goal in testing Compile/QMF was to use it as a tool for migrating QMF queries to production and to save CPU resources. "But," notes Adams, "after finding out how easy it is to use, we are going to give it to our developers for prototyping applications."

For the technical staff, the benefit is increased programmer productivity. For the end user, the benefit is increased responsiveness. "We can be more responsive to user requests," says Adams, "and we can free up our programmers to do more complex tasks."

Compile/QMF provides an easy way around one of the big drawbacks to QMF: its limitations for certain types of reports, including any reports with IF THEN logic. True, Compile/QMF will only compile the existing logic into the COBOL program. But once the COBOL code is generated, it is relatively simple to modify, adding the required additional logic. (Compile/QMF generates programs that are structured and well-commented with meaningful paragraph names.)

Reports with IF logic were the driving force for trying Compile/QMF for John Mackintosh, manager of applications development for Black Box Corporation (Pittsburgh, PA). Black Box, a manufacturer and distributor of data communications equipment, recently upgraded its order entry system. One result was a backlog of reports that had to be converted to the new system. While many of the reports have been generated in QMF, many more are too cumbersome, needing multiple selects, builds and joins in QMF to replace a single IF statement.

Compile/QMF should provide Black Box with a cost-effective solution. "We can still do the QMF query, use Compile/QMF to generate the COBOL code and add the IF statement," says Mackintosh. That was the initial reason for installing Compile/QMF. Now that it is in-house, Mackintosh sees potential incremental value for it in conjunction with an automated scheduling tool they are also installing. "The automated scheduler needs return codes from SQL errors. By converting our QMF queries to COBOL we could generate the true error codes we need."

Like QMF, however, Compile/QMF does not do it all. Compile/QMF currently cannot handle save data. Another problem for some users is the reliance on external routines for certain formatting functions.

This is because Compile/QMF was designed to accommodate those features of QMF that cannot be handled with COBOL, specifically word wrapping. By using external routines, the end user does not lose those features when the reports are converted to COBOL. (SableSoft reports that the next release, currently being field tested, does handle save data and offers the option of eliminating the external routines.)

Technical Support

Technical support, especially for new products, is always a paramount concern. SableSoft offers 24-hour technical support for Compile/QMF. Both Donahue and Adams asked for technical support and got what they needed expeditiously. Donahue reported that he found four minor problems in initial testing and they were all corrected in about a day. (The problems he found involved rounding, numbers of places to the right of the comma, a heading value not appearing on the first page of the report and failure to block the work file.)

Adams, too, was pleased with the responsiveness when he asked for help, especially considering that he was asking for help doing something the product was not designed to do. He explains, "The product is designed to be run from a TSO terminal, using TSO Attach. I wanted to run in the production control environment using Call Attach so I called SableSoft and they gave me the JCL for relinking the modules."

Regarding evaluation criteria for the three users surveyed, all of them claimed productivity and performance benefits are more important than cost, customer support or the company's background. Donahue said the cost was attractive and he expects the payback period to be short; however, he was more concerned about getting COBOL code that is maintainable. For Adams, the issue was performance and for Mackintosh, the key issue was ease of use. At \$13,000 to \$16,900 for the first license, it should not take long to cost-justify Compile/QMF.

Compile/QMF operates with any IBM 370 or plug-compatible system. It requires MVS or VM, DB2 or SQL/DS, and a COBOL compiler. For more information, contact SableSoft, Inc., 2695 Winding Trail Drive, Boulder, CO 80304, (303)443-7791. ⊜

ABOUT THE AUTHOR

Renee Peterson reports on mainframe software developments and is a frequent contributor to MAIN-FRAME JOURNAL. worst possible case is to have zero MATCHING COLUMNS, but even this may be acceptable for an index-only scan.

The usage of 100 pages as the cutoff value for the index space scan report may need to be modified. Some suggestions would be to use the same number as chosen for the tablespace scan report or to choose a value that is appropriate for the environment.

The pertinent point to be made is to carefully analyze the results of these queries. The SQL for each query is listed in the report. Every query listed on the report is not necessarily a problem query. Each query, however, should be closely monitored. Corrective actions can be taken for poorly performing queries identified via these reports. Some of these actions are listed below:

- Ensure that appropriate indexes are available
- Change current indexes to include additional columns used in a problem query, which necessitates dropping and re-creating the index a potentially long process for large tables
- Recode the queries for efficiency (that is joins versus sub-selects, use stage 1 predicates, eliminate usage of OR and so on.)
- Eliminate any unnecessary joins.

Summary

Many reports can be produced from the DB2 catalog to aid in performance monitoring. This article details some sample reports in three basic areas: navigational, physical and plan. Each area provides a proactive means of preventing performance bottlenecks before they occur. They are only suggestions. Changing them to suit specific needs should be easy due to the ad hoc nature of SQL. Good luck in tuning. \triangleq

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years of experience in database management systems. He is also a vice-president and co-founder of ASSET, Inc., a Pittsburghbased consulting and software development firm.